The loop update is an example of a cluster algorithm. Detour: Cluster algorithm for the Ising model

Define bond index corresponding to pair of interacting spins

bond  $b = 1, 2, \ldots, N_b$ , interacting spins  $\sigma_{i(b)}, \sigma_{j(b)}$ 

Number of bonds  $N_b = dN$  for a d-dimensional cubic lattice Write the energy of the Ising ferromagnet as

$$E = -|J| \sum_{b=1}^{N_b} [\sigma_{i(b)}\sigma_{j(b)} + 1] = -\sum_{b=1}^{N_b} E_b$$

Write the partition function as

$$Z = \sum_{\sigma} e^{-E(\sigma)/T} = \sum_{\sigma} \prod_{b=1}^{N_b} e^{E_b/T} = \sum_{\sigma} \prod_{b=1}^{N_b} [1 + (e^{E_b/T} - 1)]$$

Define bond functions with arguments 0,1 (bond variable):

$$F_b(0) = 1$$
  

$$F_b(1) = e^{E_b/T} - 1$$
  

$$Z = \sum_{\sigma} \prod_{b=1}^{N_b} [F_b(0) + F_b(1)]$$

Introduce bond variables

AT.

$$\tau_b = 0, 1, \ \tau = \{\tau_1, \tau_2, \dots, \tau_{N_b}\}$$

Partition function can be written as sum over spins and bonds

 $\Lambda T_{-}$ 

$$Z = \sum_{\sigma} \prod_{b=1}^{N_b} [F_b(0) + F_b(1)] = \sum_{\sigma} \sum_{\tau} \prod_{b=1}^{N_b} F_b(\tau_b)$$

The functions  $F_b$  depend on the spins:

$$F_b(0) = 1$$
  

$$F_b(1) = e^{E_b/T} - 1 = \begin{cases} e^{2|J|/T} - 1, & \text{if } \sigma_{i(b)} = \sigma_{j(b)} \\ 0, & \text{if } \sigma_{i(b)} \neq \sigma_{j(b)} \end{cases}$$

 $\tau_b = 1$  allowed only between parallel spins Probabilities: For everything else fixed, probability for a given b

$$P(\tau_b) = \frac{F(\tau_b)}{F(0) + F(1)} = \frac{F(\tau_b)}{e^{2|J|/T}}$$

If parallel spins on bond b, probabilities for the bond variable

$$P(\tau_b = 0) = e^{-2|J|/T}, \quad P(\tau_b = 1) = 1 - e^{-2|J|/T}$$

If anti-parallel spins on bond b

$$P(\tau_b = 0) = 1, \quad P(\tau_b = 1) = 0$$

For a fixed bond configuration, spins forming clusters (spins connected by "filled" bonds) can be flipped and then give a configuration (term) with the same weight in Z ( $F_b=1$ for all bonds between clusters,  $F_b$  unchanged inside cluster).



 $N(\tau_b = 1) =$  No. of filled bonds  $W = (e^{2|J|/T} - 1)^{N(\tau_b = 1)}$ (unchanged after flip)

Spins not connected to any filled bonds are single-spin clusters

# Swendsen-Wang algorithm

- Start from spin configuration
- Generate bond configuration
- Identify clusters of spins connected by bonds
- Flip each cluster with probability 1/2
- Generate new bonds with the current spins, etc

# **SSE: Linked vertex storage** for loop update

The "legs" of a vertex represents the spin states before (below) and after (above) an operator has acted



5

0

2   •	3   0
0	•
	1
0	1

3

1



 links to next and previous leg

Spin states between operations are redundant; represented by links

network of linked vertices will be used for loop updates of vertices/operators

#### **Pseudocode: Sweep of loop updates**

constructing all loops, flip probability 1/2

```
do v_0 = 0 to 4L - 1 step 2

if (X(v_0) < 0) cycle

v = v_0

if (random[0 - 1] < \frac{1}{2}) then

traverse the loop; for all v in loop, set X(v) = -1

else

traverse the loop; for all v in loop, set X(v) = -2

flip the operators in the loop

endif

enddo
```

 visited vertices are no longer needed and we set them to a negative value -1 or -2, to indicate that the loop has been visited (-1) or visited and flipped (-2)

construct and flip a loop

 $v = v_0$ do X(v) = -2 p = v/4; s(p) = flipbit(s(p), 0) v' = flipbit(v, 0) v = X(v'); X(v') = -2if  $(v = v_0)$  exit enddo

- p is the location of the operator in the original length-L list of operatotors
- by flipping bit 0 of s(p), the operator changes from diagonal to offdiagonal, or vice versa
- moving on the vertex to the adjacent spin is also done with a bit flip

#### We also have to modify the stored spin state after the loop update

- we can use the information in  $V_{first}$ () and X() to determine spins to be flipped
- spins with no operators,  $V_{first}(i)=-1$ , flipped with probability 1/2

do 
$$i = 1$$
 to  $N$   
 $v = V_{\text{first}}(i)$   
if  $(v = -1)$  then  
if  $(\text{random}[0-1] < 1/2) \sigma(i) = -\sigma(i)$   
else  
if  $(X(v) = -2) \sigma(i) = -\sigma(i)$   
endif  
enddo

v=V<sub>first</sub>(i) is the location of the first vertex leg on site i

- flip the spin if X(v)=-2
- (do not flip it if X(v)=-1)
- no operation on i if  $v_{first}(i)=-1$ ; then it is flipped with probability 1/2

### **Constructing the linked vertex list**

Traverse operator list *s(p)*, *p*=0,...,*L*-1 • vertex legs **v=4p,4p+1,4p+2,4p+3** 

Use arrays to keep track of the first and last (previous) vertex leg on a given spin

- V<sub>first</sub>(i) = location v of first leg on site i
- V<sub>last</sub>(i) = location v of last (currently) leg
- these are used to create the links
- initialize all elements to -1

$$\begin{array}{l} V_{\rm first}(:) = -1; \ V_{\rm last}(:) = -1 \\ {\rm do} \ p = 0 \ {\rm to} \ L - 1 \\ {\rm if} \ (s(p) = 0) \ {\rm cycle} \\ v_0 = 4p; \ b = s(p)/2; \ s_1 = i(b); \ s_2 = j(b) \\ v_1 = V_{\rm last}(s_1); \ v_2 = V_{\rm last}(s_2) \\ {\rm if} \ (v_1 \neq -1) \ {\rm then} \ X(v_1) = v_0; \ X(v_0) = v_1 \ {\rm else} \ V_{\rm first}(s_1) = v_0 \ {\rm endif} \\ {\rm if} \ (v_2 \neq -1) \ {\rm then} \ X(v_2) = v_0; \ X(v_0) = v_2 \ {\rm else} \ V_{\rm first}(s_2) = v_0 + 1 \ {\rm endif} \\ V_{\rm last}(s_1) = v_0 + 2; \ V_{\rm last}(s_2) = v_0 + 3 \\ {\rm enddo} \end{array}$$

creating the last links across the "time" boundary

do i = 1 to N  $f = V_{\text{first}}(i)$ if  $(f \neq -1)$  then  $l = V_{\text{last}}(i)$ ; X(f) = l; X(l) = f endif enddo

### **Determination of the cut-off L**

- adjust during equilibration
- start with arbitrary (small) n

Keep track of number of operators n

- increase L if n is close to current L
- e.g., *L=n+n/3*

### Example

- 16×16 system,  $\beta$ =16  $\Rightarrow$
- evolution of L
- n distribution after equilibration
- truncation is no approximation



## Does it work? Compare with exact results

- 4×4 exact diagonalization
- Bethe Ansatz; long chains

## Susceptibility of the 4×4 lattice $\Rightarrow$ $\approx$

- SSE results from 10<sup>10</sup> sweeps
- improved estimator gives smaller error bars at high T (where the number of loops is larger)





## ⇐ Energy for long 1D chains

- SSE results for 10<sup>6</sup> sweeps
- Bethe Ansatz ground state E/N
- SSE can achieve the ground state limit (T→0)

# **Improved Estimators**

SSE with loop updates is an example of a cluster algorithms - we can utilize improved estimators for many observables

**Classical example: Swendsen-Wang Ising cluster algorithm** 



 $N(\tau_b = 1) =$  No. of filled bonds  $W = (e^{2|J|/T} - 1)^{N(\tau_b = 1)}$ (unchanged after flip)

Write magnetization as sum over clusters of size  $n_c$ , sign  $s_c$ :

$$M = \sum_{i=1}^{N} \sigma_i = \sum_{C=1}^{N_{\rm clus}} \sum_{i \in C} \sigma_i = \sum_{C=1}^{N_{\rm clus}} s_C n_C \qquad \langle M^2 \rangle = \sum_{C=1}^{N_{\rm clus}} \sum_{C'=1}^{N_{\rm clus}} \langle n_C n_{C'} s_C s_{C'} \rangle$$

All cluster orientations (signs) have same weight - average over all 2<sup>Nclus</sup> orientations →

$$\langle M^2 \rangle = \sum_{C=1}^{N_{\rm clus}} \langle n_C^2 \rangle$$

This is the improved estimator of <M<sup>2</sup>> - only depends on cluster structure

# **Improved estimators in SSE**

Consider a given slice (propagated state) of an SSE configuration

 label the sites according to the loops passing through → clusters

In given loop All spins on given sub-lattice A or B are same, different on A, B



Staggered magnetization on a cluster is 1/2 of the size of the cluster

- changes sign when loop flipped
- similar to magnetization in SW algorithm

$$\langle M_{z,{\rm stagg}}^2\rangle = \frac{1}{4}\sum_{C=1}^{N_{\rm clus}} \langle n_C^2\rangle$$

The uniform magnetization requires the staggered phases

$$\chi = \frac{\beta}{4N} \left\langle \sum_{j=1}^{C} \left( \sum_{i=1}^{n_j} \phi_i \right)^2 \right\rangle \qquad \phi_i = \begin{cases} +1 & i \text{ on A site} \\ -1 & i \text{ on B site} \end{cases}$$

# Transverse-field Ising model [PRE 68, 056701]

$$H = \sum_{i,j} J_{ij} \sigma_i^z \sigma_j^z - h \sum_i \sigma_i^z$$

Arbitrary interactions (incl. random, long-range,...) Elementary operators for the SSE strings:

$$H_{0,0}=1,$$

$$H_{i,0}=h(\sigma_i^++\sigma_i^-), \quad i>0,$$

$$H_{i,i}=h, \quad i>0,$$

$$H_{i,j}=|J_{ij}|-J_{ij}\sigma_i^z\sigma_j^z, \quad i,j>0, \ i\neq j$$
(a)  $\bigcirc 0$ 
(b)  $\bigcirc 0$ 
(c)  $\bigcirc 0$ 
(c)  $\bigcirc 0$ 
(d)  $\bigcirc 0$ 
(d)  $\bigcirc 0$ 
(e)  $\bigcirc 0$ 
(e)  $\bigcirc 0$ 
(f)  $\bigcirc$ 

$$H = -\sum_{i=1}^{N} \sum_{j=0}^{N} H_{i,j} \qquad S_n = [i(1), j(1)], \dots, [i(n), j(n)]$$
$$Z = \frac{1}{L!} \sum_{\alpha} \sum_{S_L} \beta^n (L-n)! \langle \alpha | \prod_{l=1}^{L} H_{i(l), j(l)} | \alpha \rangle$$

Possible local updates:

$$[0,0]_{p} \leftrightarrow [i,j]_{p}, \quad i,j \neq 0, \qquad \text{(diagonal)}$$
$$[i,i]_{p_{1}}[i,i]_{p_{2}} \leftrightarrow [i,0]_{p_{1}}[i,0]_{p_{2}}, \quad i \neq 0 \qquad \text{(off-diagonal)}$$



# **Cluster update for Transverse-field Ising models**

Start at arbitrary in-leg and apply weight-preserving operator changes

- growing "to do" list (stack) from each out-leg
- until stack is empty
- divides the entire system into clusters
- similar to classical S-W algorithm

# **Long-range interactions**

Goal: avoid summations over interactions J<sub>ij</sub>

- we can first assume that all diagonal operators are allowed, then

$$P([0,0] \rightarrow [i,j]) = \frac{\beta \left( Nh + 2\sum_{ij} |J_{ij}| \right)}{L - n + \beta \left( Nh + 2\sum_{ij} |J_{ij}| \right)} \qquad P([i,j] \rightarrow [0,0]) = \frac{L - n + 1}{L - n + 1 + \beta \left( Nh + 2\sum_{ij} |J_{ij}| \right)}$$

where [i,j] in [0,0]->[i,j] is a still undetermined interaction

- the actual [i,j] is chosen in a second step using cumulative probabilities:

$$P_{c}(k) = \frac{\sum_{i=1}^{N} P(i)}{\sum_{i=1}^{N} P(i)}$$

N is the total number of diagonal ops Choose op #k using bisection in the table  $P_c()$ - time scales as ln(N)

Cluster update remains the same for any kind of interaction

- total time for updating sweep scales as βNIn(N) with long-range interactions

